

# A Monte-Carlo study of meanders

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**Abstract.** We study the statistics of meanders, *i.e.* configurations of a road crossing a river through  $n$  bridges, and possibly winding around the source, as a toy model for compact folding of polymers. We introduce a Monte-Carlo method which allows us to simulate large meanders up to  $n = 400$ . By performing large  $n$  extrapolations, we give asymptotic estimates of the connectivity per bridge  $R = 3.5018(3)$ , the configuration exponent  $\gamma = 2.056(10)$ , the winding exponent  $\nu = 0.518(2)$  and other quantities describing the shape of meanders.

**PACS.** 64.60.-i General studies of phase transitions – 05.10.Ln Monte Carlo methods – 02.10.Eb Combinatorics

## 1 Introduction

The concept of folding has an important place in polymer physics [1,2]. Considering a polymer chain made of  $n$  identical constituents (the *monomers*), the entropy of such a system can be obtained by counting the number of inequivalent ways of folding the chain onto itself. If the model of polymer does not take self-avoidance into account, it is then equivalent to the well-known Brownian motion. Several more involved models have been proposed which study the *compact* folding of a *self-avoiding* polymer as a Hamiltonian cycle (*i.e.* a closed, self-avoiding walk which visits each vertex) on a regular lattice [3,4]. They can also be defined for some kinds of *random* lattices [5,6], where each configuration is now described by a system of non-intersecting arches which connect the pairs of monomers, which are neighbors in the real space.

In the present paper, the compact folding of a polymer chain is modeled by a folded strip of stamps, with a complete piling of the strip on top of one stamp [7]. It is then equivalent to another model of non-intersecting arches, the so-called *meander* problem, which can be summarized by this simple question: in how many ways  $M_n$  can a road cross a river through  $n$  bridges, and possibly wind around the source.

A related problem can be defined by forbidding the winding around the source (*i.e.* the river is infinite at the both ends). It is now equivalent to enumerate the “simple alternating transit mazes” [8] of depth  $n$ ; it was also investigated in connection with Hilbert 16<sup>th</sup> problem, namely the enumeration of ovals of planar algebraic curves [9].

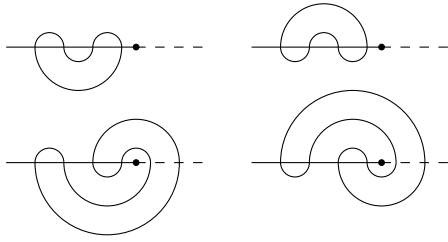
By analogy with some models of statistical mechanics like random walks or self-avoiding walks, the meanders can be described in the language of critical phenomena.

In particular, the asymptotic behavior of meanders when  $n$  is large can be characterized by “critical” exponents. However the exact enumeration of meanders is particularly complicated: there is no known formula for  $M_n$  in terms of  $n$ . By generating all possible configurations, by hand or with a computer, the beginning of the sequence  $M_n$  can be computed [10–15] exactly. As  $M_n$  increases exponentially with  $n$ , the limits of computers are reached for  $n \sim 30$  and the estimates of the exponents are too inaccurate to validate (or invalidate) some conjectures.

One should mention that several exact results, for arbitrarily large  $n$ , which are unfortunately not helpful to determine the values of the exponents, have been obtained with other techniques: random matrix model methods [13,16–18] and an algebraic approach using the Temperley-Lieb algebra [19,20] or the Hecke algebra [21].

Many models in statistical physics can be studied by Monte-Carlo (or stochastic) methods. With these algorithms, only a small set of configurations among all the possible ones are generated. In principle, the expectation of physical quantities (like energy or magnetization) with a given law of probability (like the Boltzmann law involving an external temperature) can be estimated from these randomly generated samples, if their probabilities of generation are known. For example, with the Metropolis algorithm for classical spin systems [22], the probability of generation is built to be equal to the Boltzmann law and the average is done over the generated configurations with equal weights. To bypass some difficulties (for example when the phase space has many local minima with high free energy barriers between them), it is possible, in principal, to generate the random configurations with another more adapted law and to correct this bias when the average is done [23].

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**Fig. 1.** The  $M_4 = 4$  meanders of size 4. The road is the non-self-intersecting loop. The semi-infinite river is the solid half-line, starting at the source (black dot). The size is the number of bridges. The winding number  $w$  is the number of arches crossing the dashed half-line on the right side. The upper meanders have no winding ( $w = 0$ ), but the lower have  $w = 2$ .

But for the meanders problem, the situation is quite different: the phase space is not easy to built because the number  $M_n$  of configurations is unknown and the only known efficient method to draw a meander of size  $n$  is a recurrence over  $n$ . Moreover the naïve way to use this recurrence gives a distribution of meanders which is not flat, and this default increases exponentially with  $n$ . This paper presents a Monte-Carlo method which explores the meanders with an almost flat distribution law. Furthermore the bias is known and can be corrected exactly. Therefore, the average can be done over meanders with equal probabilities. In particular, better estimates of critical exponents are obtained.

After Section 2 devoted to the definitions, and Section 3 which explains the building of meanders by recurrence, Section 4 of this paper describes the Monte-Carlo method. The results are presented and discussed in Section 5.

## 2 Definitions

A meander of size  $n$  is a planar configuration of a non-self-intersecting loop (road) crossing a half line (semi-infinite river with a source) through  $n$  points (bridges). Two meanders are considered as *equivalent* if their roads can be continuously deformed into each other, keeping the bridges fixed: this is therefore a topological equivalence. We call an arch each section of road between two consecutive bridges. So a meander of size  $n$  has  $n$  bridges and  $n$  arches.

The number of different meanders of size  $n$  is denoted by  $M_n$ . For example,  $M_1 = 1$ ,  $M_2 = 1$ ,  $M_3 = 2$ ,  $M_4 = 4$ . In Figure 1, the 4 meanders of size 4 are drawn. The  $M_n$ 's, up to  $n = 29$ , can be found in references [14,15].

In previous articles [13,14,19], these objects were called semi-meanders, to distinguish them from the case where the line is infinite (river without source). In this paper, the river is always a half-line and the word meanders is used for convenience.

We can define [13,14,19] meanders with  $k$  connected components, *i.e.* made of one river and  $k$  non-intersecting roads. But, in this work, we do not include this generalization and we keep  $k = 1$ . However the Monte-Carlo

method, used in this article, can be adapted without difficulties to an arbitrary fixed  $k$ , and even for varying  $k$  with a fugacity  $q^k$ .

As explained with many details in reference [14], the meander problem is absolutely equivalent to the problem of the compact folding of a strip of stamps because each meander of size  $n$  can be continuously deformed in such a way that the “road” becomes a vertical line and the “river” becomes a folded strip of  $n - 1$  stamps. We prefer to present our results with the meander representation because the main recursion relation, described later, seems more “natural” in this picture.

The meander problem has certain similarities with two-dimensional self-avoiding walks: a meander is obtained by intersecting a closed self-avoiding walk by a half-line and keeping only the topological aspect. By analogy, it is expected [16] that

$$M_n \stackrel{n \rightarrow \infty}{\sim} c \frac{R^n}{n^\gamma}, \quad (1)$$

where the estimates given in reference [14] are  $R = 3.50(1)$  and  $\gamma \simeq 2$ .

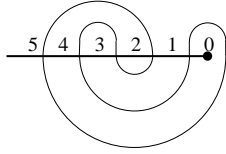
The connectivity  $R$  can be reinterpreted as the average number of ways of adding a bridge close to the source by deforming an arch of a given large meander. Then  $\ln(R)$  is the entropy per bridge. The configuration exponent  $\gamma$  is sensitive to the boundary conditions, for example whether the road is closed or open, whether the river is infinite or semi-infinite, straight or forked, whether the meander is drawn on a planar surface, a sphere or a surface with higher genus. Conversely, we expect that  $R$  remains the same for all these boundary conditions.

It is similar in on-lattice self-avoiding walks problem where the connectivity depends on the type of lattice (square, honeycomb ...) and not on the boundary conditions, whereas the “universal” configuration exponent depends on the boundary conditions, but is not sensitive to the small scale details of the lattice. For these reasons, we think that the numerical value of  $R$  is valid only for this particular model of meanders. But  $\gamma$  is expected to be more “universal” and to keep its value in other variants of the meander problem. Unfortunately, the numerical determination of  $\gamma$  is less precise than  $R$ , because  $n^\gamma$  describes the correction to the leading exponential asymptotic behavior  $R^n$ .

For a given meander  $m$ , the winding number  $w(m)$  of the road around the source of the river can be defined as the minimal number of intersections between the road and a half (semi-infinite) line starting at the source and extending the river on the opposite side. For an example, (see Fig. 1). We define  $w_n$  as the average of the winding number

$$w_n = \frac{1}{M_n} \sum_{m=1}^{M_n} w(m) \quad (2)$$

over all the meanders  $m$  of size  $n$ . We can see the winding number as the topological end-to-end distance between the source (right end of the river) and the infinite (left



**Fig. 2.** The height  $h(i)$  (resp.  $h(-i)$ ) is the number of arches over (resp. below) the segment  $i$ . For this meander of size 5,  $\{h(i)\} = \{0, 1, 2, 3, 2, 1, 0, 1, 2, 1, 0\}$  for  $i = -5, \dots, 5$ . The area is  $A = 13$ .

end of the river). Here the distance between two points is simply the minimal number of arches which must be crossed to go from one point to the other. By analogy with the end-to-end exponent of self-avoiding walk, we expect that

$$w_n \stackrel{n \rightarrow \infty}{\sim} n^\nu, \tag{3}$$

where the estimate given in reference [14] is  $\nu = 0.52(1)$ .

If we study the meanders by leaving free the number  $k$  of connected components, the problem is equivalent [14] to a random walk on a semi-infinite line and can be studied with usual methods of combinatorics. In particular,  $\gamma = 3/2$ ,  $R = 4$  and  $\nu = 1/2$  is the Brownian exponent. But, by fixing  $k = 1$ , the problem is drastically more difficult and the above values are, to our knowledge, not yet known exactly.

For a given meander  $m$  of size  $n$ , we label by  $i = 0, \dots, n$  each segment of river in-between two consecutive bridges, from right to left. So the rightmost segment (with source) is labeled 0, and the leftmost (semi-infinite) segment is labeled  $n$ . We define the height  $h(i, m)$  as the number of arches over the segment  $i$ , and  $h(-i, m)$  as the number of arches below the segment  $i$ . An example is given in Figure 2.

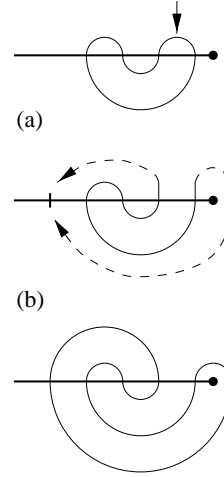
For the case  $i = 0$ , the both definitions  $h(+0, m)$  and  $h(-0, m)$  are equivalent and equal to the winding number:  $h(0, m) = w(m)$ . From the definition, we have  $h(n, m) = h(-n, m) = 0$ ,  $h(i, m) \geq 0$  and  $h(i + 1, m) = h(i, m) \pm 1$ .

For a given meander  $m$  of size  $n$ , we define the area  $A(m)$  as

$$A(m) = \sum_{i=-n}^n h(i, m). \tag{4}$$

For meanders of size  $n$ , it can be proved that the maximal area is  $(n - 1)^2 + 1$  for the two meanders with a snail shape (where  $\{h(i)\} = 0, 1, 2, \dots, n - 2, n - 1, n - 2, \dots, 2, 1, 0, 1, 0\}$  plus the symmetric meander), and the minimal area is  $2n - 2$  (resp.  $2n - 1$ ) when  $n$  is even (resp. odd) for the  $2^{n/2-1}$  (resp.  $2^{(n-1)/2}$ ) snake shaped meanders characterized by  $h(i) + h(-i) = 2$  for  $0 < i < n$ . As in the case of the winding number, we will consider the average profile height

$$h_n(i) = \frac{1}{M_n} \sum_{m=1}^{M_n} h(i, m) \tag{5}$$



**Fig. 3.** A meander of size  $n+1$  is built from a meander of size  $n$  with one labeled exterior arch by the following process. (a) Add a new bridge on the left side of the river. Cut the labeled arch. Stretch its two free ends. (b) Close the arch on the opposite side by crossing the new bridge (possibly by bypassing the source on the right). The inverse process is the following. (b) Open the road at the place of the left most bridge. (a) Pull the two free ends and close them on the opposite side to form a exterior arch.

and the average area

$$A_n = \frac{1}{M_n} \sum_{m=1}^{M_n} A(m) \tag{6}$$

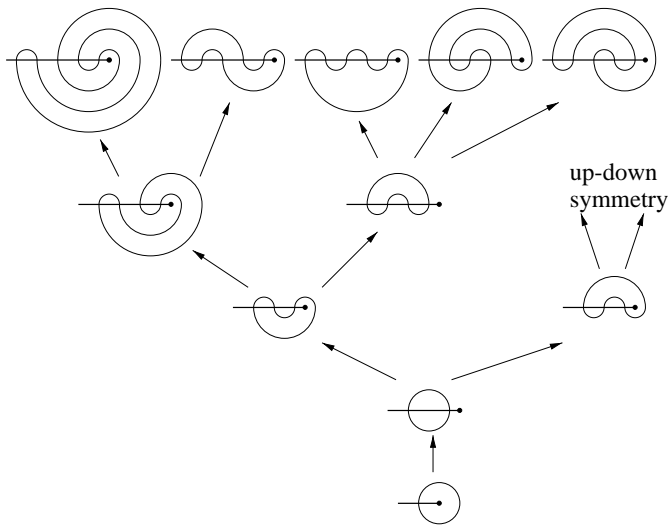
over all the meanders  $m$  of size  $n$ .

### 3 Recursion relation for meanders

In this section, we describe a recursive algorithm to enumerate and built all meanders of a given size  $n$ . Though it was described in references [13,14], we prefer to recall it in details, because the Monte-Carlo method is based on this recursion.

We have different ways to built a meander of size  $n+1$ , starting from a meander of size  $n$ . Our method consists in adding a bridge on the left most part of the river (opposite to the source) and changing the road to cross this new bridge. To keep this change minimal, only an exterior arch is modified (an arch is exterior when no other arch surrounds it).

Take a meander of size  $n$  (the parent) and choose (or label) one of its exterior arch. By the process described in Figure 3, a meander of size  $n + 1$  (the child) is built. A parent has as many different children as exterior arches. By inverting this process, it appears that each meander of size  $n + 1$  has one and only one parent. More precisely, it is a one-to-one mapping between the meanders of size



**Fig. 4.** The tree of meanders up to  $n = 5$ . For  $n \geq 4$ , only a half of the tree is drawn by using the up-down symmetry. Then  $M_1 = 1$ ,  $M_2 = 1$ ,  $M_3 = 2$ ,  $M_4 = 4$  and  $M_5 = 10$ . Each arrow represents a process as described in Figure 3.

$n + 1$ , and the meanders of size  $n$  with one labeled exterior arch.

The starting point of the recursion is the unique meander of size 1. By  $n - 1$  successive applications of the recursion process, every meanders of size  $n$  can be built. As shown in Figure 4, the set of meanders is organized as a tree. The root, at level 1, is the starting meander  $n = 1$ . Each branch between a node on level  $n$  and a node on level  $n + 1$  represents a relation between a parent of size  $n$  and its child of size  $n + 1$ . Apart from  $n = 1$ , a meander (or node) has several exterior arches, then several children (or branches). Their number depends on the precise shape of the parent and varies between 2 and  $n/2 + 1$ .

If we want to exactly enumerate the  $M_n$  meanders of size  $n$ , the only method we know is to built and investigate the tree up to the level  $n$ . In particular, we have not found a direct recursion between the numbers  $M_n$ . The number of children of each meander has a distribution which seems to be erratic and the only way to know is the examination of its shape. Then, to compute  $M_n$ , the work is proportional to  $M_n$ . As the  $M_n$ 's increase exponentially, the limits of the capabilities of the computers are rapidly reached.

In references [14,15], the meanders numbers  $M_n$  up to  $n = 29$  are given. With a recent computer and more tricks of programming, it is perhaps possible to obtain  $n = 31$  or 32. If the power of computers continues to increase exponentially, the best we can expect with the full enumeration without major improvement, is a linear growth in  $n$ , with a rate of only one new size every 2 or 3 years. We have the feeling that this progress is too small to change significantly our understanding of the meanders problem.

In this article, we will investigate larger  $n$  with a Monte-Carlo method increasing more slowly than an

exponential. But, this stochastic method gives results with error bars and the interpretation is more delicate.

## 4 The Monte-Carlo method

As explained in the previous section, the exponential growth of the computations with an exact enumeration method limits the size of meanders around  $n \sim 30$ . To reach bigger  $n$ , it is then natural to try to study this problem with a Monte-Carlo (or stochastic) method.

As the set of the  $M_n$  meanders of a given size  $n$  is too large to be fully explored, the general idea is to randomly select a small subset. Then, the measurements are done and averaged on the selected meanders. It gives an estimator of the exact (but unknown) result, with a unknown error. This error has two components: statistical fluctuations and bias.

The statistical fluctuations can be reduced by independently repeating the procedure many times. Then, we obtain a histogram of the estimator, with an average and a variance. Under the hypothesis of finite variance, the statistical fluctuations of the average can be estimated by usual formulas of statistics.

The bias is the difference between the exact result and the mathematical expectation of the estimator. If it can be exactly calculated, we subtract it from the estimator. But, in general, an unknown part remains, which can not be reduced by a better statistics. As explained below, by adjusting parameters of the simulation, the bias can be reduced to become smaller than the statistical fluctuations.

In this section, we will first introduce the simplest algorithm, the one-squirrel method. We will see that its statistical fluctuations grow exponentially with  $n$  and they are too big for  $n \approx 30$ . Then, we present an algorithm, the multi-squirrel method, for which the fluctuations increase less rapidly.

### 4.1 The one-squirrel method

The method is based on the recursion relation (see Fig. 3), with which the set of meanders is organized as a tree (see Fig. 4). The Monte-Carlo squirrel has the following stochastic behavior. It starts at the root of the tree (the meander of size  $n = 1$ ). It climbs into the tree. At each level  $n$ , it stands on a node and makes some measurements concerning the meander of size  $n$ , represented by this node. Then the squirrel goes to the level  $n + 1$  by choosing at random one of the  $b_n$  branches, starting from this node. The squirrel stops at a prefixed level  $n = n_{\max}$ . This process constitutes one simulation.

The probability that the squirrel reaches a given meander of size  $n$  is  $1/\prod_{l < n} b_l$ . This probability law is not flat because the sequence of  $b_l$  depends on the visited nodes. As seen in Figure 4, for  $n = 5$ , the two meanders on the left side have a probability  $1/8$  and the three on the right side have  $1/12$ . So, to correct this bias between meanders,

the squirrel has a weight

$$q_n = \prod_{l=1}^{n-1} b_l, \quad (7)$$

which is calculated during its climbing.

By noting  $\langle \cdot \rangle$  the mathematical expectation (over all possible simulations), it is then obvious that  $\langle q_n \rangle = M_n$ , because the sum runs over  $M_n$  possible paths and the contribution of a given path is  $q_n$  with probability  $1/q_n$ .

More generally, for some quantity  $Z$  (for example the winding number), we wish to determine

$$\mathcal{Z} = \sum_{m=1}^{M_n} Z(m), \quad (8)$$

where  $Z(m)$  is the value of  $Z$  on the  $m$ -th meander of size  $n$ . Each simulation gives, at level  $n$ , a measurement  $Z(s)$  on the meander  $s$  reached by the squirrel and  $\langle q_n Z(s) \rangle = \mathcal{Z}$  which is the generalization of  $\langle q_n \rangle = M_n$  obtained with  $Z = 1$ . Then

$$z = q_n Z(s) \quad (9)$$

is an unbiased estimator of  $\mathcal{Z}$  (*i.e.*  $\langle z \rangle = \mathcal{Z}$ ).

As usual in Monte-Carlo methods, several simulations are made independently and we hope that the average  $\bar{z}$  of all the measurements  $z$  is close to  $\mathcal{Z}$ . Unfortunately, this method does not work in practice, because the weight  $q_n$  is the product of  $b_l$ . Although the distribution of each  $b_l$  is regular, the product of many random variables is not self-averaging.

As the sum of  $\ln(b_l)$  is self-averaging (*i.e.* the observed result is closed to its mathematical expectation when  $n$  goes to infinite), most of the observed  $q_n$  are not closed to  $\langle q_n \rangle$  and

$$\frac{\langle q_n \rangle}{q_n(\text{observed})} \sim \exp \sum_{l < n} (\ln \langle b_l \rangle - \ln b_l) \quad (10)$$

increases like an exponential. Then the averages with weight  $q_n$  are dominated by exponentially rare events and the statistical fluctuations become large. To keep the observed average close to the mathematical expectation, the number of simulations must increase exponentially with  $n$  and fluctuations become too big for  $n \sim 30$  or  $35$ . As the difficulties increase exponentially with  $n$  (as for exact enumeration), it is useless to increase the power of the computer. We need a new algorithm.

## 4.2 Multi-squirrels method

We generalize the one-squirrel method, but now with a population of  $S$  squirrels, which reproduce and die;  $S$  is a fixed parameter during all the simulations. It is more simple to choose  $S$  as  $S = M_{n_0}$ , with at the starting point, a squirrel staying at each node of level  $n_0$  (meanders of size

$n_0$ ). In this work,  $n_0 = 17$  and we use the up-down symmetry to reduce the population to  $S = M_{17}/2 = 1664094$  squirrels.

The population evolves from level  $n$  to level  $n + 1$  by the following process. Each squirrel  $i$  lives on a node  $s_i$  on level  $n$ , connected to  $b_i$  nodes on level  $n + 1$ . It reproduces and has  $b_i$  children and each child lives on each one of these  $b_i$  nodes. The total number of children  $S' = \sum_{i=1}^S b_i$  is calculated. The ratio  $B_n = S'/S$  is an estimate of  $M_{n+1}/M_n$ : the average of the number of children per parent. To prevent an exponential growth of the population and of the needed computer memory and time, the total population is kept constant by decimating the children: only  $S$  among the  $S'$  children survive. The choice is made at random with uniform distribution. Then the probability of surviving is  $1/B_n$ . This decimation is the single Monte-Carlo step of the algorithm.

This process is iterated up to reach a prefixed level  $n = n_{\max}$ : it gives one simulation. Many independent simulations are done and averaged.

The particular case  $S = 1$  gives the previous method with one squirrel. But, as for  $S = 1$ , for every value of  $S$ , the probability that a given meander is reached, is not uniform. The nodes with small number of ‘‘brothers’’ or ‘‘cousins’’ have always a small advantage. But this bias becomes smaller when the population  $S$  is large. That is the main improvement of this method. The limit  $S = \infty$  corresponds to the exact enumeration.

To correct the bias, each simulation has a weight

$$q_n = \prod_{l=n_0}^{n-1} B_l \quad (11)$$

and the averages runs over all the simulations with their weight. More exactly, for some quantity  $Z$ , by keeping the notations of equation (8), one simulation with  $S$  squirrels gives  $S$  measurements  $\{Z(s_i)\}$  for  $i \in [1, S]$  with a weight  $q_n$  and the estimator

$$z = q_n \sum_{i=1}^S Z(s_i) \quad (12)$$

is unbiased, *i.e.*

$$\langle z \rangle = \mathcal{Z}. \quad (13)$$

We note that the case  $Z = 1$  gives  $S \langle q_n \rangle = M_n$ .

In order to prove equation (13), we define the operator  $\delta_m$  characterizing a given meander  $m$  by  $\delta_m(m') = 1$  when  $m = m'$  and 0 otherwise. Then every operator  $Z$  can be split up into  $Z = \sum_m Z(m) \delta_m$ . As the expectation of a sum is always the sum of expectations, we have to prove equation (13) for the operators  $\delta_m$  only, which becomes  $\langle q_n \Delta_m \rangle = 1$ , where  $\Delta_m = 1$  if the  $m$ -th meander is occupied by a squirrel and 0 otherwise. Let  $p$  represent the parent of  $m$  in the tree at level  $n - 1$ . The probability that  $\Delta_m = 1$  (*i.e.*  $m$  is occupied) is the product of the probability  $\Delta_p$  that  $p$  was occupied at the level  $n - 1$  and the probability  $1/B_{n-1}$  that its child  $m$  survives after the decimation process ( $n - 1 \rightarrow n$ ).

By using equation (11) and averaging on the random decimation ( $n - 1 \rightarrow n$ ) only,

$$\langle \Delta_m \prod_{l=n_0}^{n-1} B_l \rangle = \langle \Delta_p \prod_{l=n_0}^{n-2} B_l \rangle. \quad (14)$$

It is a recursion relation between a given meander of size  $n$  and its parent of size  $n - 1$ . By iterating, we go down to the ancestor at the starting level  $n_0$  for which  $\Delta = 1$  and the empty product of  $B_l$  is 1. It proves that  $\langle q_n \Delta_m \rangle = 1$  for every meander  $m$  of size  $n$ , and equation (13) is valid for every operator  $Z$ .

As usual, many simulations are done and the measurements are averaged with their respective weight. *A priori*, it seems that this method has the same defect as the one-squirrel method because the weight  $q_n$  is the product of many  $B_l$ , not self-averaging when  $n$  becomes large. The ratio

$$\frac{\langle q_n \rangle}{q_n(\text{observed})} \sim \exp \sum_{l=n_0}^{n-1} (\ln \langle B_l \rangle - \langle \ln B_l \rangle) \quad (15)$$

between the mathematical expectation and the most frequently observed  $q_n$  increases like an exponential. But, the main improvement of the multi-squirrel method is that the distribution of  $B_l$  becomes narrow when the population  $S$  of squirrels is large. As  $B_l = S'/S$ , the fluctuations of  $B_l$  are of order  $O(1/\sqrt{S})$  because the number  $S' = \sum_i b_i$  of children is the sum of  $S$  random variables. A Taylor expansion of  $\ln B_l$  shows that  $\ln \langle B_l \rangle - \langle \ln B_l \rangle = O(1/S)$ .

Then, with these simple arguments, we can hope that the ratio (15) grows like  $1 + O(n/S)$  and that problems appear only when  $n$  becomes on the same order than  $S$ . In our simulations, we observe that the fluctuations grow with  $n$  faster than this optimistic prediction  $O(n/S)$ . In fact, the  $B_l$ 's are not independent and the exponential function accentuates all deviations. Then we supervised carefully the distribution of  $q_n$ . When  $n$  is small, we see a regular bell-shaped curve. But, when  $n$  increases, the distribution becomes asymmetric, with a long and irregular tail for the large  $q_n$ .

For example, for  $n = 400$ , with  $S = 1\,664\,094$  squirrels, the width  $\sigma$  of the distribution is only 12%, but we observed rare events with a value of  $q_n$  as big as three times the average. However, in this case, their contribution to the average and fluctuations is not yet problematic. But, if we let  $n$  increase without control, rare events will dominate and the results will become hazardous.

How to choose  $n$  and  $S$ ? The naïve point of view is to take  $n$  as bigger as possible. But, to make  $N_s$  independent simulations with  $S$  squirrels of size  $n$ , the need for computer memory is of order  $O(nS)$  and the need for computer time is of order  $O(n^2 SN_s)$ . If  $S$  is large enough, the fluctuations are Gaussian and the error bars are of order  $O(1/\sqrt{SN_s})$ . As  $n$  is always limited, we will extrapolate to study the asymptotic behavior. For that, it is of no help to have large values of  $n$  if the error bars are too big. So for a fixed computer time, we prefer accumulate good statistics by limiting  $n \leq 400$ . Finally for a fixed product  $SN_s$ , we

prefer to take  $S = 1\,664\,094$  as bigger as permitted by the memory computer to avoid the problem of rare but large fluctuations.

### 4.3 Bias for non-linear observables

In the previous section, we have seen how to obtain unbiased Monte-Carlo estimates of the sum  $\mathcal{Z}$  over all the meanders of size  $n$  of some quantity  $Z$  (see Eq. (8) and its notations). However we are more interested by the average  $\mathcal{Z}/M_n$  over all the meanders of size  $n$ . For example, the average winding number  $w_n$  (see Eq. (2)) is obtained when  $Z$  counts the winding. To evaluate  $R$  of equation (1), we can analyze  $M_{n+1}/M_n$ ; in this case,  $Z$  counts the exterior arches. More generally, we want to use non-linear combinations of  $\mathcal{Z}$  and  $M_n$ .

With our Monte-Carlo method, we have seen that one simulation gives a measurement  $z$  which is an unbiased estimator:  $\langle z \rangle = \mathcal{Z}$ . With  $N_s$  independent simulations, we call  $\bar{z}$  the usual average of the  $N_s$  measurements  $z$ ; its fluctuations are  $\sqrt{N_s}$  times smaller. The bar over the symbols distinguishes the average of observed values by Monte-Carlo method, from the (unknown) mathematical expectation, marked with  $\langle \dots \rangle$ . The same work can be done with  $q_n$  which is a unbiased estimator of  $M_n$ .

We must be careful with the Monte-Carlo estimate of  $\mathcal{Z}/M_n$ . For example, the average of the ratio  $z/q_n$  gives bad results. It is better to compute the ratio of the averages  $\bar{z}/\bar{q}_n$ . Indeed, with a Taylor expansion of  $z$  and  $q_n$  around their mathematical expectations  $\mathcal{Z}$  and  $M_n$ , the bias (defined as the difference of the mathematical expectation  $\langle \bar{z}/\bar{q}_n \rangle$  and the target  $\mathcal{Z}/M_n$ )

$$\langle \bar{z}/\bar{q}_n \rangle - \mathcal{Z}/M_n = O(1/N_s). \quad (16)$$

It can be neglected if it is smaller than the stochastic fluctuations. Usually, in Monte-Carlo simulations, this problem disappears because several millions of independent measurements are done. But, in this work, the situation is quite different. In fact, each simulation is a complex process involving millions of squirrels, and the number  $N_s$  of simulations is small.

Of course, we cannot compute this bias exactly, otherwise we would have already subtract it from measurements. But we can estimate it by the following process. The set of simulations is divided into  $N_s/2^p$  subsets, with  $2^p$  simulations each. In each subset,  $\bar{z}/\bar{q}_n$  is computed. We obtained  $N_s/2^p$  independent values, one for each subset; by usual formulae of statistics, we compute their average  $E^{(p)}$  and the error bars. This work is done for all integer  $p$  between  $p = 0$  (each subset contains only one simulation) and  $p = \log_2 N_s$  (only one set with all the  $N_s$  simulations).

Which value of  $p$  is the best? Following equation (16), the bias of  $E^{(p)}$  is expected to decrease like  $1/2^p$ . For small values of  $p$ , we observe really a dependency of  $E^{(p)}$  on  $p$ : the bias is visible. But, for  $p > 5$ , variations become smaller than statistical error bars: the size  $2^p$  of subsets is sufficiently large to neglect the bias. But, if  $p$  is close to its maximum, the number of subsets becomes very small and

the error bars are not properly estimated. As in our work  $N_s = 8192 = 2^{13}$ , we finally keep  $p = 7$ : the estimator  $E^{(7)}$  is computed with 64 independent subsets of 128 simulations each. This method was used for all quantities presented below. It is valid, not only for ratios like  $\mathcal{Z}/M_n$ , but also for non-linear functions like  $\ln(M_{n+1}/M_n)$ . It could also be possible to use more complex estimators. For example, the combination  $2E^{(p)} - E^{(p-1)}$  cancels the order  $1/2^p$  of the bias.

## 5 Results

In this section, we describe the results obtained by our Monte-Carlo multi-squirrels method. After several tests, we used a population of  $S = 1\,664\,094$  ( $= M_{17}/2$ ) squirrels for meanders with size up to  $n = 400$ . To do  $N_s = 8192$  independent simulations, we have used during 8 days a parallel computer (the Cray T3E of the CEA-Grenoble) with 128 processors (Dec-alpha at 375 MHz) and 13 Giga-bytes of total memory, equivalent to 24000 hours of workstation cpu time.

We have verified that the results are stable when  $S$  (the population) increases. More exactly, the tests with smaller  $S$  have larger error bars, but are compatible with results and error bars obtained with the largest  $S$ . As explained in the Section 4.2, we have carefully checked that  $S = 1\,664\,094$  is sufficiently large to explore sizes of meander up to  $n = 400$ .

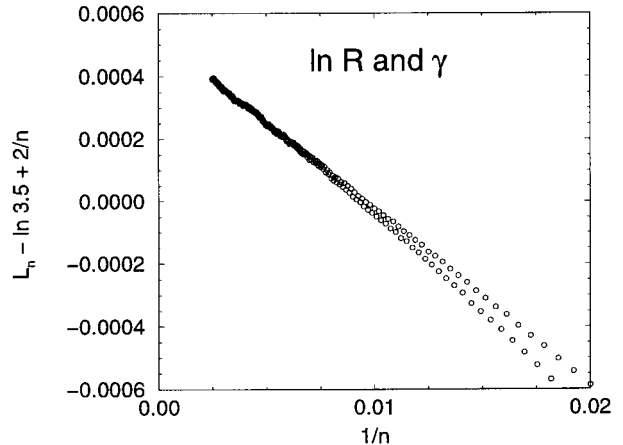
### 5.1 Enumeration

We want to measure  $R$  and  $\gamma$ , which describe the asymptotic behavior of the number of meanders  $M_n \sim c R^n/n^\gamma$  for large  $n$ . The entropy  $\ln R$  can be estimated by  $\ln(M_n/M_{n-1})$ . But it appears that the sub-sequences  $M_{2n}$  and  $M_{2n+1}$  have an alternating sub-leading correction. We have estimated it to be  $u(-1)^n/(n \ln n)$  with  $u = 0.5(1)$ . This alternating effect is dramatically amplified by the ratio  $M_n/M_{n-1}$ . So it is better to consider

$$L_n = \frac{1}{2} \ln(M_n/M_{n-2}), \quad (17)$$

with a jump from  $n-2$  to  $n$ . But even with this precaution, the reader can still see on the following figures a small parity effect. To estimate  $L_n$ , we have used the procedure described in Section 4.3.

As we expect  $L_n \sim \ln R - \gamma/n$  for large  $n$ , by plotting  $y = L_n$  versus  $x = 1/n$ , we can estimate  $\ln R$  (limit when  $x$  goes to 0) and the exponent  $\gamma$  (asymptotic slope). In Figure 5, we have plotted the Monte-Carlo estimate of  $L_n - \ln 3.5 + 2/n$  versus  $1/n$  for  $n$  between 50 and 400. We have arbitrarily subtracted the linear function  $y = \ln 3.5 - 2x$ , to reduce the amplitude of  $y$ ; we obtain a figure where the small quantities  $2 - \gamma$  (remaining slope),  $\ln(R/3.5)$  (limit when  $x$  goes to 0) and curvature (deviation to the expected linear behavior) are more visible. The curvature remains small and a linear extrapolation gives a



**Fig. 5.**  $\ln R$  and  $\gamma$ : plot of the Monte-Carlo estimate of  $L_n - \ln 3.5 + 2/n$ , for  $n$  between 50 and 400, versus  $1/n$ . The limit when  $x$  goes to 0 is  $\ln(R/3.5)$ , and the (negative) slope is  $2 - \gamma$ . The error bars are not drawn; their maximum is  $10^{-5}$ , then they are smaller than the symbols. A parity effect, between the odd and even  $n$ , is visible. A linear extrapolation gives  $R = 3.5019(2)$  and  $\gamma = 2.056(10)$ .

limit between 0.0005 and 0.0006 with an estimated slope 0.056(10). Then

$$R = 3.5019(2) \quad \text{and} \quad \gamma = 2.056(10). \quad (18)$$

With the assumption that the asymptotic behavior is  $M_n \sim c R^n/n^\gamma$ , the conjecture  $\gamma = 2$  is incompatible with these simulations. But, we can try another asymptotic shape, for example

$$M_n \sim c \frac{R^n}{n^\gamma} \frac{1}{\ln^\alpha n}, \quad (19)$$

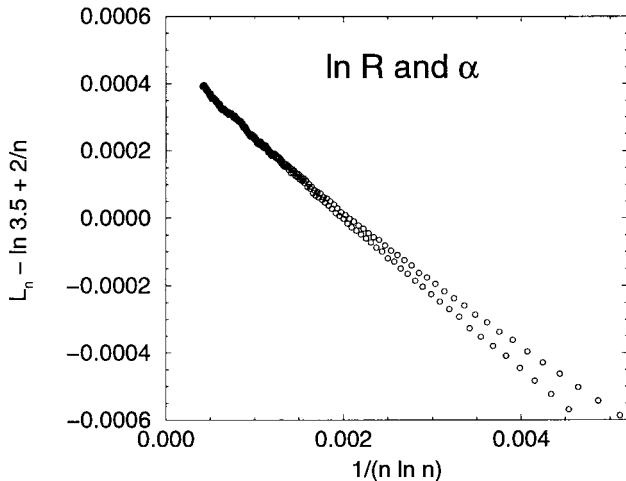
by introducing a new exponent  $\alpha$ . In Figure 6, we have plotted  $L_n - \ln 3.5 + 2/n$  (as in Fig. 5), but now versus  $1/(n \ln n)$ . With this transformation of the  $x$ -axis, a linear behavior corresponds to  $\gamma = 2$  and the slope measures  $-\alpha$ . A linear extrapolation gives

$$R = 3.5017(2), \quad \gamma = 2 \quad \text{and} \quad \alpha = 0.25(5). \quad (20)$$

with  $\alpha$  compatible with the simple fraction  $1/4$ .

How shall we choose between both results equation (18, 20)? We notice that the quality of the alignment of points is the same in Figures 5, 6. In fact, it is very difficult to distinguish between a logarithmic law and a power law with such a small exponent, with statistical error bars and when the amplitude of  $\ln n$  is small. In fact, for  $n$  close to  $n'$ , the term  $\alpha/(n \ln n)$  looks like  $\alpha'/n$  with  $\alpha'/\alpha = 1/\ln n' + 1/\ln^2 n'$ . Then, for any choice of  $\alpha$  not too large,  $\gamma = 2.056 - 0.22\alpha$  gives a class of acceptable behaviors.

We have tried to use more sophisticated extrapolation methods. For example, with a fixed jump  $i$ ,  $(nL_n - (n-i)L_{n-i})/i$  gives theoretically the same limit  $\ln R$  but by removing the term  $\gamma/n$ . Thus  $n^2(L_n - L_{n-i})/i$  gives a direct estimate of  $\gamma$ . Unfortunately these kinds of derivative



**Fig. 6.**  $\ln R$  and  $\alpha$  under the hypothesis  $\gamma = 2$ : the same plot as Figure 5 but the  $x$ -axis is  $1/(n \ln n)$ . The (negative) slope is  $-\alpha$ . A linear extrapolation gives  $R = 3.5017(2)$  and  $\alpha = 0.25(5)$ .

amplify statistical errors, and the results are compatible but less precise than the previous estimates.

So, with our numerical simulations, we cannot say if a logarithmic factor is present or not. However,  $R = 3.5018(3)$ , and we can exclude  $R = 3.5$  and the conjecture  $\gamma = 2$  without logarithmic factor ( $\alpha = 0$ ).

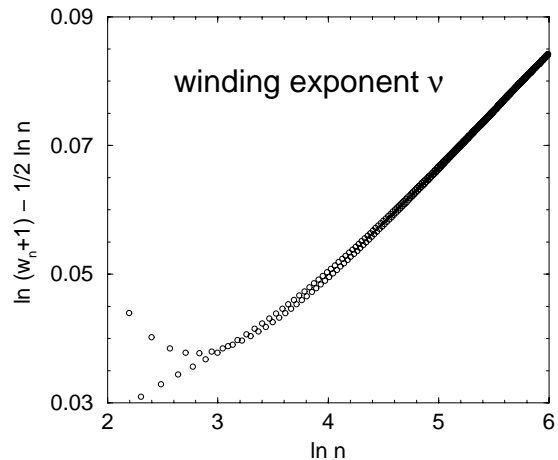
## 5.2 Winding

We will present the Monte-Carlo results of the exponent  $\nu$ , which describe the asymptotic behavior of the average winding number  $w_n \sim n^\nu$ . To avoid problem with bias, we have used the procedure described in Section 4.3. By plotting  $\ln(w_n)$  versus  $\ln(n)$ , the asymptotic slope will be a measurement of  $\nu$ . In Figure 7, we have plotted  $y = \ln(w_n + 1) - 1/2 \ln n$  versus  $x = \ln n$ . We have arbitrarily considered  $\ln(w_n + 1)$  and not  $\ln(w_n)$  because  $w_n + 1$  is less sensitive than  $w_n$  to the finite size effects [14]. As the main question is to know whether  $\nu = 1/2$  or not, we have arbitrarily subtracted the linear function  $y = x/2$ . Then the variation of  $y$  is reduced and we obtain a figure where  $\nu - 1/2$  (residual slope) and the curvature are more visible. We see that the curvature is small and a linear extrapolation gives  $\nu = 0.518$ . As it is difficult to estimate the errors with the data of Figure 7, we have also tried more sophisticated quantities like

$$G_i(n) = \frac{n}{i} \ln \left( \frac{w_n + 1}{w_{n-i} + 1} \right) \quad (21)$$

which are discrete derivatives of  $\ln(w_n + 1)$  with step  $i$ . They give a direct value for  $\nu$ , but unfortunately the statistical fluctuations are amplified by this differentiation and the uncertainty over  $\nu$  is of order 0.002.

We have seen that, for the exponent  $\gamma$ , a behavior with logarithmic correction is not excluded by the Monte-Carlo



**Fig. 7.** Winding exponent  $\nu$ : plot of the Monte-Carlo results of  $\ln(w_n + 1) - 1/2 \ln n$ , for  $n$  between 8 and 400, versus  $\ln n$ . The slope is 0.018; it is a measurement of  $\nu - 1/2$ . The error bars are not drawn; their maximum is  $3 \times 10^{-4}$ , then they are smaller than the symbols.

data. So we have tried to fit the winding number with

$$w_n \sim n^{1/2} \ln^\alpha n. \quad (22)$$

With this hypothesis, a plot of  $y = \ln(w_n + 1) - 1/2 \ln n$ , as in Figure 7, but now versus  $x = \ln \ln n$  would give a straight line with slope  $\alpha$ . But the curvature is much stronger than that of Figure 7. So we dismiss this hypothesis and conclude that

$$\nu = 0.518(2). \quad (23)$$

With the assumption that the asymptotic behavior is a simple power-law, the Brownian value  $\nu = 1/2$  is incompatible with these simulations.

## 5.3 Probability distribution of winding number

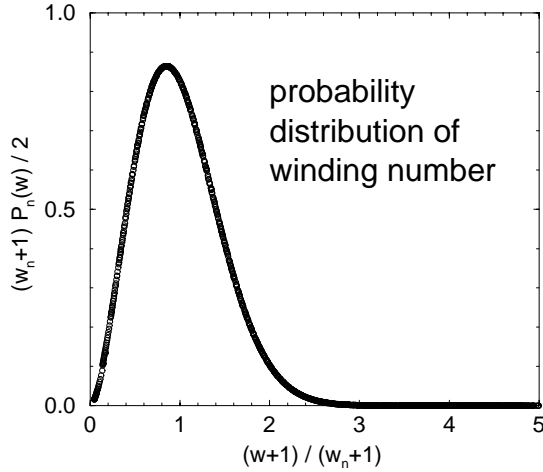
We define the probability distribution  $P_n(w)$  of winding number as the fraction of the meanders of size  $n$  with  $w$  windings. We expect [14] the asymptotic scaling behavior

$$P_n(w) \approx \frac{2}{w_n + 1} f \left( \frac{w + 1}{w_n + 1} \right) \quad (24)$$

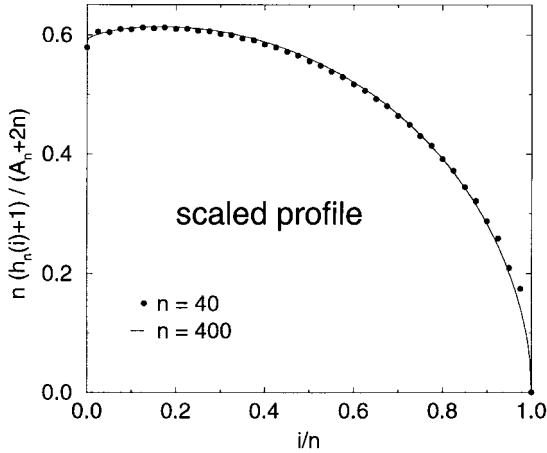
with a scaling function  $f(x)$ , where  $w_n$  is the average winding number. With the factor 2, the integral of  $f$  is normalized to 1 because  $w$  and  $n$  are integers with the same parity. In Figure 8, we plot  $y = (w_n + 1)P_n(w)/2$  versus  $x = (w + 1)/(w_n + 1)$  for different values of  $n$ . To define the scaling variable  $x$ , we prefer to take  $w + 1$  instead of  $w$  to reduce finite size effects.

We see that the points accumulate on a smooth curve, which represents the scaling function  $f(x)$ . By analogy with the end-to-end distribution for polymers, we expect [14] a power law behavior,  $f(x) \sim x^\theta$ , for small  $x$ ,





**Fig. 8.** Plot of  $(w_n + 1)P_n(w)/2$  versus  $(w + 1)/(w_n + 1)$  for  $n = 30, 40, 50 \dots 400$ . The points accumulate on a scaling function  $f$ . The error bars are not drawn; their maximum is  $6 \times 10^{-4}$ .



**Fig. 9.** Plot of the scaled height  $n(h_n(i) + 1)/(A_n + 2n)$  versus the scaled coordinate  $x = i/n$ . Only two set of data ( $n = 40$  and  $400$ ) are shown in order not to overload the figure. The points accumulate on a scaling function  $\rho(x)$ . We can see small deviations for  $x = 0$  and  $x = 1$ , which are analyzed in the text.

and a behavior  $f(x) \sim \exp(-\text{const} \times x^\delta)$ , for large  $x$ . Our data give the estimates

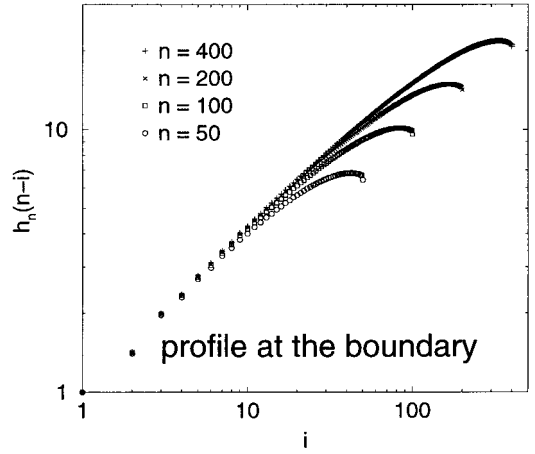
$$\theta = 1.7(1) \quad \text{and} \quad \delta = 2.3(1). \quad (25)$$

To obtain a better precision, it would require to have larger values of  $n$ .

#### 5.4 Height and area

We are interested by the asymptotic behavior, for large size  $n$ , of the average area  $A_n$  (see Eq. (6)) and average height  $h_n(i)$  (see Eq. (5)).

The label  $i$  is the ‘‘horizontal’’ coordinate and varies between  $-n$  and  $n$ . So, we introduce the scaled variable



**Fig. 10.** Log-log plot of the height  $h_n(n-i)$  versus the distance to the boundary  $i$  for various sizes  $n$ . The points accumulate on a limiting curve, which can be fitted by  $i^\phi$  with  $\phi = 0.64(2)$ .

$x = i/n$ , with  $-1 \leq x \leq 1$ . We expect [14] that

$$h_n(nx) \sim \frac{A_n}{n} \rho(x) \quad (26)$$

for large  $n$  by fixing  $x$ , where  $\rho(x)$  is a scaling function with integral normalized to 1. In Figure 9, we have plotted  $y = n(h_n(i) + 1)/(A_n + 2n)$  versus  $x = i/n$  for various size  $n$ . Only the positive  $i$  are shown because  $h_n(i)$  is symmetric after summing over all the meanders. More precisely, we have plotted the average of  $\{h_n(i) + h_n(-i)\}/2$  over the Monte-Carlo samples of meanders, which is equivalent to the average of  $h_n(i)$  over these samples, plus those obtained by the left-right symmetry ( $i \rightarrow -i$ ). As in the previous figures, we prefer to take  $(h_n + 1)$  and  $(A_n + 2n)$  to reduce finite size effects.

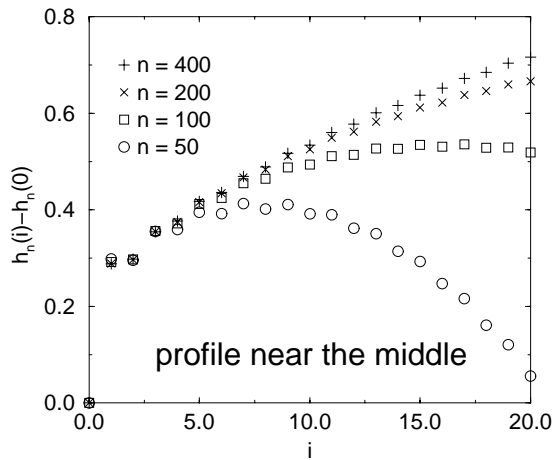
We see that, for  $0 < x < 1$ , the points accumulate on a smooth curve, which represents the scaling function  $\rho(x)$ . This shape is not a half-circle, as it would be for a random-walk on a semi-infinite line [14].

As  $h_n(0)$  is the winding number  $w_n$  for the particular case  $x = 0$ , equation (26) can be valid only if  $h_n(nx)$  scales as  $n^\nu$  with the same exponent  $\nu$  for all  $x$ . Consequently, the area  $A_n$  scales as  $n^{\nu+1}$ . In a previous section, we have numerically determined  $\nu = 0.518(2)$  by extrapolation of  $w_n = h_n(0)$ . The same work with  $h_n(n/2)$  and  $A_n$  gives  $\nu = 0.517$  for both, compatible with the previous estimates, but less precise because the finite size effects are stronger.

In Figure 9, small deviations appear between the curves for  $n = 40$  and  $400$ , at the boundary ( $x \simeq \pm 1$ ) and in the middle of the meander ( $x \simeq 0$ ).

In order to understand why finite-size effects are important near the boundary, Figure 10 is a plot of  $h_n(n-i)$  versus  $i$ , with various  $n$ . Clearly, when  $n$  is large, curves accumulate on a limiting curve

$$\tilde{h}(i) = \lim_{n \rightarrow \infty} h_n(n-i). \quad (27)$$



**Fig. 11.** Plot of the  $h_n(i) - h_n(0)$  versus  $i$  for small  $i$ , near the middle of meander, for various size  $n$ . By symmetry, only the positive  $i$  are drawn. As the heights  $h_n(i)$  are shifted by  $h_n(0)$  and not scaled, we observe effects smaller than the asymptotic behavior  $n^\nu \rho(0)$ , plus finite corrections of order  $\hat{h}(i)$ .

As this function can be fitted by a straight line on this “log-log” plot, we define a new exponent  $\phi$  by

$$\tilde{h}(i) \sim i^\phi \quad \text{with} \quad \phi = 0.64(2) \quad (28)$$

when  $i$  is large. As by construction,  $h_n(n) = 0$  and  $h_n(n-1) = 1$  for all  $n$ ,  $\tilde{h}(0) = 0$  and  $\tilde{h}(1) = 1$  satisfy exactly equation (28).

If we set  $i = ny$ , by using equation (28) valid when  $i$  is fixed, we obtain  $h_n(n-i) \sim (ny)^\phi$ . On the other hand, by using equation (26) which is valid when  $y$  is fixed, we obtain now  $n^\nu \rho(1-y)$ . As the exponents  $\nu$  and  $\phi$  differs, these two regimes are incompatible and the behavior of  $h_n(n-i)$  depends on the order in which  $n$  and  $i$  go to infinity. In particular, the domain of validity of equation (28) is reduced to the single point  $x = 1$  for equation (26). That explains why the rightmost dots ( $i \sim n$ ) in Figure 9 for the small size are not superimposed on the curve obtained for the large size.

Here, the exponent  $\phi$  is the surface critical exponent, while  $\nu$  is the bulk critical exponent. Near the boundary,  $h(i)$  is small and the condition  $h(i) \geq 0$  limits appreciably the fluctuations toward the bottom. This effect is so strong that the exponent is changed and  $\phi > \nu$ . This situation is reminiscent of other critical phenomena [24], like the self-avoiding walk near a surface. Our result is to be contrasted with the case of a random walk on a semi-infinite line for which the surface exponent keeps its Brownian value  $1/2$ .

Finite-size effects observable near the middle of meander ( $x \sim 0$ ) can be explained with Figure 11, which is a plot of  $h_n(i) - h_n(0)$  versus  $i$  for small  $i$  with various  $n$ . Clearly, when  $n$  is large, curves accumulate on a limiting curve

$$\hat{h}(i) = \lim_{n \rightarrow \infty} (h_n(i) - h_n(0)). \quad (29)$$

If we make the hypothesis that the behavior of  $\hat{h}(i)$  is compatible with equation (26) by inverting the limits  $n$  large

and  $x$  small, the consequences would be that  $\rho(x)$  has a cusp at  $x = 0$  with a infinite derivative  $\rho(x) \sim \rho(0) + x^\nu$  when  $x$  is small, and  $\hat{h}(i) \sim i^\nu$ . But this power law behavior of  $\hat{h}(i)$  is not observed. Then the asymptotic behavior of  $h_n(i)$  with  $i$  fixed and  $n$  large is given by equation (26), *i.e.*  $n^\nu \rho(0)$ , plus finite corrections of order  $\hat{h}(i)$ .

This cusp is an another boundary effect since the point  $i = 0$  is the source of the river. Let us consider three consecutive heights  $\{h(i-1, m), h(i, m), h(i+1, m)\}$  for a given meander  $m$ . By definition,  $h(i+1, m) = h(i, m) \pm 1$ . Then, for a generic  $i$ , the couple of its neighbors can have 4 respective values  $\{h(i-1), h(i+1)\} = \{h(i) \pm 1, h(i) \pm 1\}$ . But, for the special case  $i = 0$ , the situation  $h(-1) = h(1) = h(0) - 1$  happens only if a single arch connects the first bridge (near the source) to itself, by drawing a little circle around the source, without visiting the other bridges. This is forbidden if we insist in having only one connected component. In other word, the neighborhood of the source limits the fluctuations of  $h(0)$  toward the top. In particular, for every meander,  $h(0) \leq \{h(-1) + h(1)\}/2$ . That explains why, on average,  $h(0) < h(1)$ .

To understand why  $h(1) < h(2)$ , it is mandatory to consider more complex forbidden situations for  $\{h(-2), h(-1), h(0), h(1), h(2)\}$ . For example, a circle connecting bridges 1 and 2 with an upper and a lower arch, with  $h(\pm 1) = h(\pm 2) + 1$ , is forbidden.

More generally, for a given  $i$ , the presence of the source forbids the systems of arches connecting the bridges  $\{1, 2, \dots, i\}$  with a closed road without visiting the other bridges  $\{i+1, \dots\}$ . Qualitatively, it is a repulsive “force” which favors the connection of bridge  $i$  with bridges  $j > i$  and gives a concave shape for small  $i$ .

As the forbidden situations are more and more complex when  $i$  becomes large, their statistical effects decrease and this repulsive force has a finite range. In the end, the summation over all forbidden situations gives finally this cusp with a finite amplitude described by  $\hat{h}(i)$ .

## 6 Conclusion

In this paper, we have presented a Monte-Carlo method to investigate a phase space described by a deterministic but irregular tree (*i.e.* the number of branches at each node is not fixed). With a naïve random climbing on the tree (the one-squirrel method), the probability of a path depends on the number of branches encounter at each node. For the meanders problem, the ratio between extremal weights increases exponentially with the size: consequently the most part of the computer time is devoted to generate configurations with small weight, and only a exponentially small number of configurations with high weight contribute efficiently to the average.

With the multi-squirrel method, the distribution becomes almost flat: the bias, *i.e.* the ratio between extremal weights increases very slowly and never exceeds 3 in our simulations. Moreover, this bias is exactly known during the simulation, then it can be corrected to average over all meanders with a uniform distribution. As usual

with Monte-Carlo simulations, results suffer from statistical fluctuations which decrease, in the best case, like the square root of the computer time.

After a simulation on a parallel computer with 3 years of cpu time in single processor units, results with small errors bars have been obtained for meanders up to size  $n = 400$ . Under some hypothesis inspired by the analogy with random walks problems, large  $n$  extrapolation can be done for the enumeration (see Eq. (18) and Fig. 5), for the distribution (see Eq. (25) and Fig. 8) and the average of the winding number (see Eq. (23) and Fig. 7) and the shape (see Sect. 5.4) of meanders.

From a Monte-Carlo point of view, the proposed algorithm can be used, in principal, for any combinatorial problem described by a tree: the essential ingredient is to know, for a given node, the number of branches. However algorithms are rare in this kind of problems and better variants or other algorithms could be without doubt invented.

About the meanders, with these estimates, it appears that the critical exponent are not simple fractions as  $1/2$  or  $7/2$ , as conjectured by previous studies [16]. Of course, Monte-Carlo simulations cannot determine the exact values, but can confirm or invalidate analytical proposals, while waiting for a rigorous solution.

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